

AEG: Automatic Exploit Generation

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How do I pwn?

- Take a binary
- Find a vulnerability and inputs which trigger that vulnerability
- Create a payload which exploits the vulnerability
- ???
- profit responsible disclosure
- Vendor doesn't patch it after several months
- profit

Why can't I pwn?

- Vulnerability discovery is a slow and tedious process
- Large size of binaries
- Vulnerability → Exploit can be nontrivial
 - e.g. restrictions on input, insufficient space for shellcode, etc.
- Patching of vulnerabilities varies in difficulty





DARPA Cyber Grand Challenge (2014-)2016

- Automatic exploitation and patching
- Custom pwnables written for DECREE OS
- DECREE caveats and rant
- Winners:
 - 1st: Mayhem (CMU)
 - 2nd: Xandra (TECHx)
 - 3rd: Mechanical Phish (UCSB)
- Only Mechanical Phish (angr) opensourced :(

Mayhem in CGC

Challenges modeled after real exploits

- Morris Worm (buffer overflow)
- Stuxnet LNK (off by one) (CVE-2010-2568)
- Crackaddr (buffer overflow) (CVE-2002-1337)
- Heartbleed (leak of sensitive data) (CVE-2014-0160)

Patching

- Return pointer encryption
- Protection of indirect calls/jmps
- Extended malloc allocations
- Manual ASLR
- Cleaning of uninitialized space

 For DEFCON challenge that was broadcasted at 14:10:25 UTC, hardened binary created at 14:11:08 UTC (43 seconds)



Symbolic Execution Primer



Toy Program

```
void foo(int x, int y){
    int t = 0;
    if (x > y) {
        t = x;
    } else{
         t = y;
    if (t < x){</pre>
        assert false;
```

```
void foo(int x, int y){ • Concrete Store:
    int t = 0;
                            -x = 4
    if (x > y) {
       t = x;
                            -y = 4
    } else{
        t = y;
    }
    if (t < x) {
        assert false;
    }
```

```
void foo(int x, int y){ • Concrete Store:
    int t = 0;
                             -x = 4
    if (x > y) {
       t = x;
                             -y = 4
    } else{
                             -t = 0
        t = y;
    }
    if (t < x){</pre>
        assert false;
    }
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```
void foo(int x, int y){ • Concrete Store:
    int t = 0;
                             -x = 4
    if (x > y) {
       t = x;
                             -y = 4
    } else{
                             -t = 4
        t = y;
    }
    if (t < x){</pre>
        assert false;
    }
```

```
void foo(int x, int y){ • Concrete Store:
    int t = 0;
                             -x = 4
    if (x > y) {
        t = x;
                             -y = 4
    } else{
                             -t = 4
        t = y;
    }
                          • 4 < 4 is false,
                            quik maffs
    if (t < x){</pre>
        assert false;
    ł
```

```
void foo(int x, int y){ • Concrete Store:
    int t = 0;
                             -x = 4
    if (x > y) {
        t = x;
                             -y = 4
    } else{
                             -t = 4
        t = y;
    }

    4 < 4 is false,</li>

                             assertion
    if (t < x){</pre>
                             unreached
        assert false;
    ł
```

```
void foo(int x, int y){ • Symbolic Store:
    int t = 0;
                                - \mathbf{X} = \mathbf{X}
    if (x > y) {
        t = x;
                                -y = Y
     } else{
         t = y;
     }
    if (t < x) {
         assert false;
     }
```

```
void foo(int x, int y){ • Symbolic Store:
    int t = 0;
                                 - \mathbf{X} = \mathbf{X}
    if (x > y) {
        t = x;
                                 -y = Y
     } else{
                                 -t = 0
         t = y;
     }
    if (t < x){</pre>
         assert false;
     }
```

}

```
void foo(int x, int y){
    int t = 0;
    if (x > y){
        t = x;
        t = x;
    }
        else{
        t = y;
    }
    if (t < x){
        assert false;
    }
    · viet of the set o
```

```
void foo(int x, int y){ • Symbolic Store:
    int t = 0;
    if (x > y) {
        t = x;
    } else{
        t = y;
    }
    if (t < x){</pre>
        assert false;
    }
```

- - $\mathbf{X} = \mathbf{X}$

- -t = ite(X < Y, X, Y)
- Assert condition: ite(X<Y,X,Y)<X</pre>

```
void foo(int x, int y){
    int t = 0;
    if (x > y) {
         t = x;
     } else{
         t = y;
     }
    if (t < x) {</pre>
         assert false;
```

- Symbolic Store:
 - $\times = X$

- -t = ite(X < Y, X, Y)
- Assert condition: ite(X<Y,X,Y)<X
- Throw into solver
 assert not hit

Dynamic Symbolic Execution

```
void foo(int x, int y){ • Symbolic Store:
    int t = 0;
    if (x > y) {
        t = x;
    } else{
        t = y;
    }
    if (t < x){</pre>
        assert false;
```

- - $\mathbf{X} = \mathbf{X}$

$$- T = 0$$

 Case split on conditional

Dynamic Symbolic Execution

```
void foo(int x, int y){ • Branch 1: X > Y
    int t = 0;
    if (x > y) {
        t = x;
    } else{
        t = y;
    }
```

}

- Symbolic Store:
 - $\mathbf{X} = \mathbf{X}$
 - -y = Y
 - t = X

- if (t < x){</pre> assert false;
- Assert condition: X<X
- Assert not hit

Dynamic Symbolic Execution

```
void foo(int x, int y){
    int t = 0;
    if (x > y){
        t = x;
    } else{
        t = y;
    }
    if (t < x){</pre>
```

```
assert false;
```

}

- Branch 2: !(X > Y)
- Symbolic Store:
 - $\times = X$

$$-y = Y$$

- -t = Y
- Assert condition:
 Y<Y
- Assert not hit



Actually Exploiting Stuff (kindof maybe)



AEG in Four Easy Steps

- Symbolically execute program (warning! slow!)
- Detect violation of safety property
- Check if exploitable
- Generate exploit (using template shellcode)



Case Study: Crackaddr Variant

• CVE2002-1337

- Sendmail 5.79 to 8.12.7
- Remote execution via buffer overflow in 'crackaddr' function of headers.c

• CGC Challenge (Halvar Flake (2011))

- Extracted core of bug (50 LOC vs. 247)
- 'Tool should automatically show vulnerable version has a bug and the fixed version is safe'

Case Study: Crackaddr Variant

```
#define BUFFERSIZE 200
 1
 2
    #define TRUE 1
   #define FALSE 0
 3
 4
    int copy_it (char *input, unsigned int length) {
 5
        char c, localbuf[BUFFERSIZE];
 6
        unsigned int upperlimit = BUFFERSIZE - 10;
7
        unsigned int quotation = roundquote = FALSE;
 8
        unsigned int inputIndex = outputIndex = 0;
 9
        while (inputIndex < length) {</pre>
10
            c = input[inputIndex++];
11
            if ((c == '<') && (!quotation)) {</pre>
12
                 quotation = TRUE; upperlimit--;
            }
13
14
            if ((c == '>') && (guotation)) {
15
                 quotation = FALSE; upperlimit++;
16
            }
17
            if ((c == '(') && (!quotation) && !roundquote) {
18
                 roundquote = TRUE; upperlimit--; // decrementation was missing in bug
19
            }
20
            if ((c == ')') && (!quotation) && roundquote) {
21
                 roundquote = FALSE; upperlimit++;
22
            }
23
            // If there is sufficient space in the buffer, write the character.
24
            if (outputIndex < upperlimit) {</pre>
25
                 localbuf[outputIndex] = c;
26
                outputIndex++:
27
            }
28
        }
29
        if (roundquote) {
30
            localbuf[outputIndex] = ')'; outputIndex++; }
31
        if (quotation) {
32
            localbuf[outputIndex] = '>'; outputIndex++; }
33
```

its a state machine woaw



- 201 loop iterations to trigger bug
- 10 different paths through loop
- 5²⁰¹ (approx 2⁶⁶⁴) paths

Case Study: Unintended Solution

1 2

3

4

5

6 7

8

9 10

11 12

13

14 15

16

17

18 19

20 21

22

23

24

26

27

```
int copy( fileSystemType *fs, char *cmdline, unsigned int owner ) {
  char sourcefile[FILENAME SIZE];
 char destfile[FILENAME_SIZE];
 int x:
 // skip over leading whitespace characters
 while ( *cmdline != 0 && isspace(*cmdline) )
   ++cmdline:
 // if we hit the end of the line there were no filenames specified
 if ( *cmdline == 0 ) {
    return ERROR BAD PARMS:
 }
 x = 0:
 while ( *cmdline != 0 && !isspace(*cmdline) ) {
   if ( x < FILENAME_SIZE ) {</pre>
      sourcefile[x] = *cmdline;
    }
   ++cmdline:
   ++x:
 -}
 sourcefile[x] = 0:
  . . .
```

Case Study: Unintended Solution

```
int main(void) {
 1
 2
      char command[1024];
 3
     while (1) {
 4
        bzero(command, 1024);
 5
        getline(command, 1024);
 6
 7
        i = 0:
 8
        while (command[i] != ' ' && i < strlen(command)) {</pre>
          ++i:
 9
10
        }
11
        command[i] = 0;
12
        if ( strcmp(command, "list") == 0 ) {
13
14
          . . .
15
        3
        else if ( strcmp( command, "copy" ) == 0 ) {
16
          retcode = copy( currentFS, command+i+1, currentUser );
18
        }
19
        . . .
20
      }
21
   }
```

Solved by Mayhem (~1h 45m)

oh no

- Symbolic execution suffers from scaling issues
- Real world nuisances like libraries, device drivers, operating systems
 - On top of standard binary analysis issues (e.g. CFG recovery)
- A lot of effort has gone into making symbolic execution of programs more viable

help! i'm too slow!

Handling path explosion

- Heuristic preconditions on state space
 - Known Length (automatic max)
 - Known Prefix (manual, e.g. HTTP GET)
 - Concolic Execution (manual, crashing input)
- Heuristic path prioritization
 - Buggy-path-first
 - Loop Exhaustion

help! i'm too slow!

Handling state space explosion

- 'Driller' architecture (Mechanical Phish)
 - Dynamic Symbolic Execution with fuzzing
 - Each shores up weaknesses of the other
- Veritesting (CMU Cylab)
 - Alternate between dynamic and static symbolic execution
 - Balances between the solver and the symbolic execution engine

help! the real world exists!

Handling the real world

- Actually symbolically execute into kernel/library
 - (probably going to fail)
- Function/syscall hooking
 - Unconstrained symbolic values
 - Model effects of function call on symbolic state
 - Tedious and possibly error prone

help! the real world exists!

Handling the real world

- Indirect jumps/calls
 - Resolve all jump targets
 - Randomly concretize
- S2E framework ('in-vivo' execution)
 - Switch between concrete and symbolic execution
 - Concretize e.g. syscall inputs, make symbolic after return

Some Remarks

- AEG is a relatively new and developing field
- Techniques have been around for decades
- Practical implementations of AEG are still very much in development
- Real world is hard
- Formal methods is (are?) cool

Useful Readings

- Symbolic Execution Survey
 - https://github.com/season-lab/survey-symbolic-e xecution

Decision Procedures, SMT solving

- The Calculus of Computation (Bradley, Manna)
- Logic in Computer Science (Huth)
- Theorem Proving/Provers
 - CPDT (Chlipala), DeepSpec project
 - CompCert, seL4
 - Coq, Isabelle/HOL, Twelf, Idris, etc. etc.



COOL VIDEO



Cool video

D:\Documents\AEG Exploits Demo.mp4



thanken you

